

by a user or subscriber of a wireless communication system. Examples of user equipment include cellular phones, wireless laptops and pagers.

The formatting of information refers to the arrangement of units of information (typically bits) into groups or blocks of information where the size of each block is defined. Suppose, for example, that the 3GPP standard uses eight modes to do its formatting. Each of the eight modes of the 3GPP standard arranges information to be transmitted into three blocks respectively referred to as Class A, Class B and Class C blocks. In UMTS, the Class A, Class B and Class C blocks are referred to as Transport Channels (TrCh) and each Transport Channel has a format called a Transport Format (TF). The modes typically are the operating modes of Codecs which are devices and/or functions within the system equipment and/or subscriber equipment that perform the coding operations. The following is a table showing the format for each of the eight modes for UMTS systems that comply with the 3GPP standard:

Mode	Class A	Class B	Class C
mode 1	81 bits	103 bits	60 bits
mode 2	65 bits	99 bits	40 bits
mode 3	75 bits	84 bits	0 bits
mode 4	61 bits	87 bits	0 bits
mode 5	58 bits	76 bits	0 bits
mode 6	55 bits	63 bits	0 bits
mode 7	49 bits	54 bits	0 bits
mode 8	39 bits	56 bits	0 bits

Table 1

After the information to be transmitted is arranged, i.e., formatted, as per one of the above modes, each of the groups of information is coded for error correction and/or error detection. Coding is a technique whereby redundancies are introduced in

information to be transmitted to protect the information from having errors due to the information having been propagated through a communication channel. Error correction coding is used to correct errors and error detection coding is used to detect errors. The formatted information is subjected to various levels of coding and information padding. In particular, for UMTS systems, the formatted information is applied to a Cyclic Redundancy Coder (CRC) and tail bits are then added to pad the coded information. The formatted, coded and padded information is then applied to a convolutional coder. The output of the convolutional coder is the information coded through special mapping of every unit of the information. For example, for a convolution coder which maps 2 bits for every 1 bit of information (i.e., Rate $\frac{1}{2}$ coding), the total number of bits for each transport channel is doubled.

Convolutional coded information for each of the three transport channels is thus generated. Convolutional coding is one type of error correction coding. CRC coding is one type of error detection coding. The information of the three channels are then multiplexed prior to transmission.

FIG. 1 depicts an example of a format of the transport channels of a UMTS that complies with the 3GPP standard. Information block 100 represents the format for a transport channel which contains A bits of information with an appendage of 8 CRC bits and 8 tail bits; this transport channel is commonly referred to as TrCh1. Information block 102 represents the format for a transport channel which contains B bits of information with 8 tail bits appended; this transport channel is commonly referred to as TrCh2. Information block 104 represents the format for a transport channel which contains C bits of information and 8 tail bits; this transport channel is commonly referred to as TrCh3.

The multiplexed information is then transmitted over a communication channel or channels of the UMTS system. In UMTS systems the information is transmitted in synchronization with a timing period called a TTI (Transmission Time Interval). Transmitting equipment and receiving equipment of the system are synchronized to the TTI. Each TTI period has a beginning and an end; three blocks as per Table 1 are transmitted during a TTI. At a receiving equipment, a Class A block

is first received followed by a Class B block and then a Class C block. The size of each of these received blocks is dictated by the mode in which the system is currently operating.

5 At a receiving equipment, the information is decoded by applying procedures that are the reverse of the procedures applied by the transmitting equipment. As with the transmitting equipment, the receiving equipment can be system or subscriber equipment. In order to properly decode the formatted information, however, the receiving equipment uses the actual coded block of information and information about
10 the formatting (i.e., size of the blocks of information over each of the transport channels) of the received block. In particular, a received coded and formatted block of information is decoded by using the actual block and the size of that block to decode (e.g., CRC decoding, convolutional decoding) the received information. For example, referring to Table 1, a Class A block of information is properly decoded
15 when that block of information is applied to a decoder of the receiving equipment and the decoding uses the correct size value (i.e., the value of 81) to perform the decoding operation. If the size value used by the decoder is incorrect, the received block will not be decoded correctly.

20 To address the issue of knowing the correct block size, the current UMTS standard uses a signaling channel within which TFCI (Transport Format Combination Indicator) information is transmitted to the receiving equipment. The TFCI contains the value that represents the size of the received block. In the example discussed above, the TFCI would contain the value of 81 for a Class A block, 103 for a Class B
25 block and 60 for a Class C block. Because the decoding of the received information is dependent upon a correct size value for a received block of information, the TFCI is typically heavily coded and made more robust so as to better handle channel anomalies thus reducing the likelihood of incurring errors. As a result, more bandwidth and power is needed to transmit the TFCI. Furthermore, even though the
30 heavy coding of the TFCI reduces the likelihood of errors occurring in the TFCI, an error bound still exists. The error bound is the best error rate that can be expected from TFCI for the type and amount of coding applied to the TFCI.

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What is therefore needed is a way of detecting the format of received information without having to use TFCI.

Summary of the Invention

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The present invention provides a method of detecting the format of received information without the use of TFCI information. The present invention provides a method for Blind Transport Format Detection (BTFD). One of the communication channels through which information is received is identified as a guiding channel of the communication system. A lookup table (or other mapping technique) is provided which contains the format of the communication channels of the system as defined by the standard being followed by the communication system. Information received over the guiding channel is extracted and is applied to decoding equipment along with guiding channel format information obtained from the lookup table. The guiding channel format information is information size value that defines the size of information blocks conveyed over a transport channel. Each guiding channel format information in the lookup table is separately used with the extracted information to perform decoding operations. The guiding channel format information that, when used to decode the extracted guiding channel information yields a correct decode, is detected as the correct format. The associated formats of the other channels as defined in the lookup table are then determined from the detected correct format of the guiding channel. Thus, the format of the received information is detected without the use of TFCI information.

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Brief Description of the Drawings

FIG. 1 depicts the format of the transport channels for a UMTS that complies with the 3GPP standard.

FIG. 2 is a flowchart depicting the method of the present invention.

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FIG. 3 is a flowchart representing an algorithm used to determine the correct format from one of the steps of FIG. 2.

Detailed Description

The present invention provides a method of detecting the format of received
 5 information without the use of TFCI information. The present invention provides a
 method for Blind Transport Format Detection (BTFD). One of the communication
 channels through which information is received is identified as a guiding channel of
 the communication system. A lookup table (or other mapping technique) is provided
 which contains the format of the communication channels of the system as defined by
 10 the standard being followed by the communication system. Information received over
 the guiding channel is extracted and is applied to decoding equipment along with
 guiding channel format information obtained from the lookup table. The guiding
 channel format information is information size value that defines the size of
 information blocks conveyed over a transport channel. Each guiding channel format
 15 information in the lookup table is separately used with the extracted information to
 perform decoding operations. The guiding channel format information that, when
 used to decode the extracted guiding channel information yields a correct decode, is
 detected as the correct format. The associated information size values of the other
 channels as defined in the lookup table are then determined from the detected correct
 20 information size value of the guiding channel. Thus, the format of the received
 information is detected without the use of TFCI information.

For ease of explanation, the method of the present invention will be explained
 in the context of a UMTS communication system that complies with the 3GPP
 25 standard with a format as per Table 1. The information is thus represented in digital
 form with the unit of information being a bit. It will be readily understood that the
 method of the present invention is applicable to other communication systems
 (wireline or wireless) that use some type of formatting scheme prior to transmitting
 information over their communication channels. Information to be transmitted are
 30 multiplexed into three different channels (called transport channels) each of which has
 a format defined as per Table 1. For purposes of discussion and ease of explanation,
 transport channel 1 (TrCh1), through which Class A bits are conveyed, is selected as
 the guiding channel. The guiding channel is one of the transport channels of the

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transport channels. The information is received with radio, processing and other well known equipment typically used to detect and receive communication signals.

In step 202, information received over the guiding channel (i.e., TrCh1) is extracted. The information received over the other transport channels are also extracted. The information both from the guiding channel and the other channels are temporarily stored in any well known memory circuit for further processing. The beginning of the received information is determined from the TTI timing with which the received equipment and transmitting equipment are synchronized. The extraction of information from the channels comprise various steps necessary to receive--synchronously with the TTI--radio signals over the various channels, demodulate said signals, convert said signals to information bits and demultiplex the information bits into separate information blocks for each of the transport channels.

In step 204 decoding operations are performed on the information extracted from the guiding channel. Each decoding operation comprises an error correction decode followed by a tail bit test and an error detection decode operation. For example, when the error correction decode is convolutional decoding and the error detection decode is CRC decoding, the following procedure is performed for each decoding operation: the information extracted from the guiding channel is applied to a convolutional decoder; the convolutional decoder, which uses a certain number of shift operations or other timing operations dictated by the information size value, outputs a block of bits containing information bits (i.e., class A bits) with appended CRC bits and tail bits; tail bits are stripped off the result of the convolutional decoding and a tail bit test is performed on the tail bits; the remainder of the information is then applied to a CRC decoder which yields a CRC pass or CRC fail result. The decoding operation as described above is performed M times where M is an integer that represents the total number of modes used for the guiding channel or the total number of information size values defined for the guiding channel. Therefore, M=8 for a system which complies with the channel format defined by Table 1. It should be noted that the method of the present invention is applicable to other values for M and is not limited to a value of M=8.

As per the 3GPP standard, the 8 tail bits originally included in the transmitted information are all "0" bits. The tail bit test yields two values: one value is the number of "1" bits occurring in the tail bits and the other value indicates whether the tail bit test failed or passed. To determine whether the tail bit test passed or failed, a threshold is arbitrarily defined (by the service provider, for example) as a certain number of "1" bits occurring in the tail bits. For example, if the threshold is set to 2, the tail bit test will be declared as passing if the number of "1" bits is 2 or less. If the number of "1" bits is 3 or more, then the tail bit test is declared as failing.

For each of the M convolutional decoding operations, after the tail bits are stripped off the output of the convolutional coder, the remaining bits (i.e., Class A bits with appended CRC bits) are applied to a CRC decoder. It should be noted that a different information size value is used in each of the M convolutional decoding operations. A CRC decode operation is therefore similarly performed M times with the remaining bits where a different information size value is used to perform each of the M distinct CRC decode operations. The result of each of the CRC decode operations, which is either a "0" for 'CRC fail' or a "1" for 'CRC pass', is stored. Therefore, at the end of the decoding operations, M sets of data are stored where each set contains the following: (a) number of "1" tail bits; (b) Tail bit test "failed" or "passed"; and (c) CRC decode "pass" or "fail". As discussed above, the CRC decode operation and the tail bit test are performed on the result of the convolutional decode operation and these operations are performed M times.

In step 206, the M sets of data resulting from the decoding operations are applied to an algorithm that decides whether any of the results indicates a correct decode. The algorithm uses results (comprising three variables) to decide whether there was a correct decode, and if so, which of the M decode operations yielded the correct decode. The three variables are the result of the CRC decode (or CRC test), the result of the tail bit test and the number of non-zero bits (or "1" bits) occurring in the tail bit. In particular, the number of "1" tail bits for the i^{th} decoding operation (or i^{th} mode), where $i = 1, 2, \dots, M$ is represented as T_i ; i is any integer equal to M or less. A tail bit test "passed" if $T_i \leq T_0$ where T_0 represents the number of "1" bits

occurring in the tail bits. T_0 is a value that can be set by the service provider and/or the manufacturer of transmitting and/or receiving equipment. T_0 is also modifiable.

K_i is a binary variable that indicates whether the tail bit test passed or failed for the i^{th} mode; that is when the tail bit test fails, K_i is set to 0 and when the tail bit test passes,

5 K_i is set to 1. C_i is a binary variable that indicates whether the CRC decode operation passed or failed for the i^{th} decode operation (or i^{th} mode) ; that is when a CRC decode operation passes, C_i is set to 1 and when a CRC decode operation fails, C_i is set to 0. Therefore the three resulting values at the end of the decoding operations on information extracted from the guiding channel are (C_i, K_i, T_i) .

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Referring momentarily to FIG. 3, one particular algorithm used for step 206 of FIG. 2 is now discussed. It should be understood that the algorithm shown in FIG. 3 is one specific procedure used to decide which one of M decode operation yields a correct decode based on the CRC decode, convolutional decode and tail bit test. It should further be understood that the method of the present invention is not limited to this one particular algorithm; other algorithms and/or techniques that use the results of decoding operations to decide whether there is a correct decode are well within the scope of the method of the present invention.

20 In step 302, the results (C_i, K_i, T_i) for each of the M decoding operations are stored. Each result set is a candidate for the correct decode. Also a value for T_0 is established. If none of the M decoding operations results yielded neither a CRC pass nor a tail bit test pass (steps 304 \rightarrow 328 \rightarrow 338), a BTFD failure is declared. In cases of BTFD failure, the receiving equipment will most likely try to determine the format from another transport channel or request that the transmit equipment re-transmit the information. If only one candidate yielded a CRC pass, a BTFD correct decode is declared for that operation (steps 304 \rightarrow 306 \rightarrow 318) where $C_i = 1$. If none of the M decoding operations yielded a CRC pass and only one yielded a tail bit test pass, that particular mode is selected as a correct decode (steps 304 \rightarrow 328 \rightarrow 330 \rightarrow 340) where $C_i = 0$ and $K_i = 1$.

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If more than one candidate passed the CRC test and only one of these candidates passed the tail bit test, that candidate is selected as the correct decode (steps 304 → 306 → 308 → 310 → 356). When more than one candidate passed the CRC test and more than one of those candidates passed the tail bit test, the only candidate out of these candidates which had a number of "1" bits less than the set threshold, T_0 , is declared the correct decode (steps 304 → 306 → 308 → 310 → 312 → 314 → 326) where $C_i=1$, $K_i=1$ and $T_i < T_0$. From step 314 if, out of the candidates who passed both the CRC test and the tail bit test, there were more than one in which $T_i < T_0$, then a BTFD failure is declared (steps 304 → 306 → 308 → 310 → 312 → 314 → 316). From step 312, when more than one candidate passed both the CRC and the tail bit tests and none of these candidates satisfy the condition $T_i < T_0$, nor the condition $T_i = T_0$, then a BTFD failure is declared (steps 304 → 306 → 308 → 310 → 312 → 358 → 364). From step 358, if only one of the candidates who pass the CRC test, the tail bit test and did not satisfy the condition $T_i < T_0$, but satisfied the condition $T_i = T_0$, that candidate is declared as the correct decode (steps 304 → 306 → 308 → 310 → 312 → 358 → 360 → 366) where $C_i=1$, $K_i=1$ and $T_i = T_0$. From step 358, if more than one candidate satisfied the condition $T_i = T_0$, then a BTFD failure is declared (steps 304 → 306 → 308 → 310 → 312 → 358 → 360 → 362).

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A BTFD failure is declared when one of the following results exists for more than one candidate:

- (a) $T_i < T_0$, $C_i=0$, and $K_i=1$ (steps 304 → 328 → 330 → 332 → 334 → 336);
- (b) $T_i = T_0$, $C_i=0$, and $K_i=1$ (steps 304 → 328 → 330 → 332 → 344 → 346 → 348)
- (c) $T_i = T_0 + 1$, $C_i=1$, and $K_i=1$ (steps 304 → 306 → 308 → 320 → 322 → 354).

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Also, a BTFD failure is declared when none of the candidates passed the CRC test and more than one of the candidates passed the tail bit test, but none of those candidates satisfy either the condition of $T_i < T_0$ or $T_i = T_0$ (steps 304 → 328 → 330 → 332 → 344 → 342).

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Finally, a correct decode is declared when the following results occur for only one of the M sets of data generated from the M decoding operations:

- (a) $T_i = T_0$, $C_i = 0$, and $K_i = 1$ (steps 304 \rightarrow 328 \rightarrow 330 \rightarrow 332 \rightarrow 344 \rightarrow 346 \rightarrow 350);
- (b) $T_i < T_0$, $C_i = 0$, and $K_i = 1$ (steps 304 \rightarrow 328 \rightarrow 330 \rightarrow 332 \rightarrow 334 \rightarrow 352);
- 5 (d) $T_i = T_0 + 1$, $C_i = 1$, and $K_i = 1$ (steps 304 \rightarrow 306 \rightarrow 308 \rightarrow 320 \rightarrow 322 \rightarrow 324).

Referring back to FIG. 2, the method of the present invention moves to step 208 after having applied an algorithm such as the one discussed above (and depicted in FIG. 3) in step 206 whereby the algorithm yielded a BTFD failure. When a correct
 10 decode has been decided from one of the M sets of decoding data, the method of the present invention moves to step 210. In step 210, the corresponding information size value of the transport channels are determined from a table lookup based on the information size value of the guiding channel that yielded a correct decode. For example, if the information size value of 81 yielded a correct decode for information
 15 received over the guiding channel (i.e., TrCh1), then according to the lookup table, the corresponding information size values for TrCh2 is 103 bits and the information size value for TrCh3 is 60 bits. The format of the transport channel is thus detected. The information size values for TrCh2 and TrCh3 can thus be used to decode the Class B and Class C information bits respectively. It should be noted that mapping techniques
 20 other than the use of a lookup table can be used. In other words, the detected information size value for the guiding channel can be used to point to the information size values of the other channels of the communication system yielding the format of the received information. Such other techniques as using the memory address or storage address of the detected information size value for the guiding channel to point
 25 to the corresponding memory addresses of the information size value of the other transport channels can be used for example.

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